

Softening the Network: Virtualization's Final Frontier

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Abstractions We've Seen

virtual memory

virtual disk volumes

virtual machines

→ the *illusion* of a thing

abstraction

no re-programming

sometimes is disruptive

VM-1

VM-3

VM-2

VM-4

PM-1

VM-1

VM-3

VM-2

PM-2

meets needs

provisioning

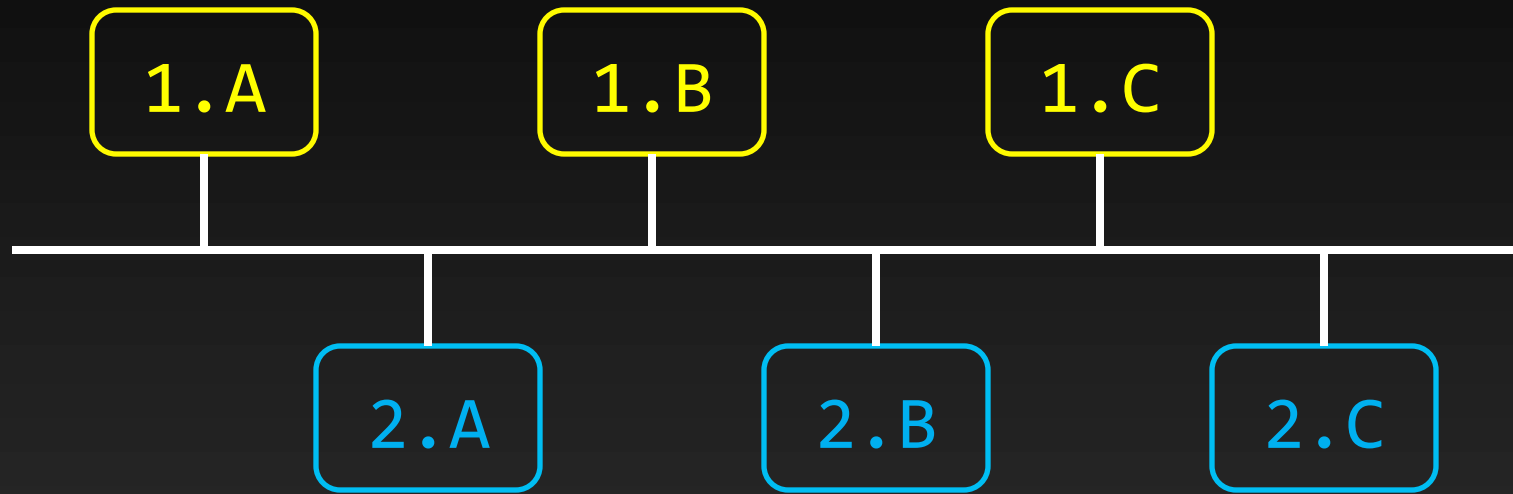
moving

snapshotting

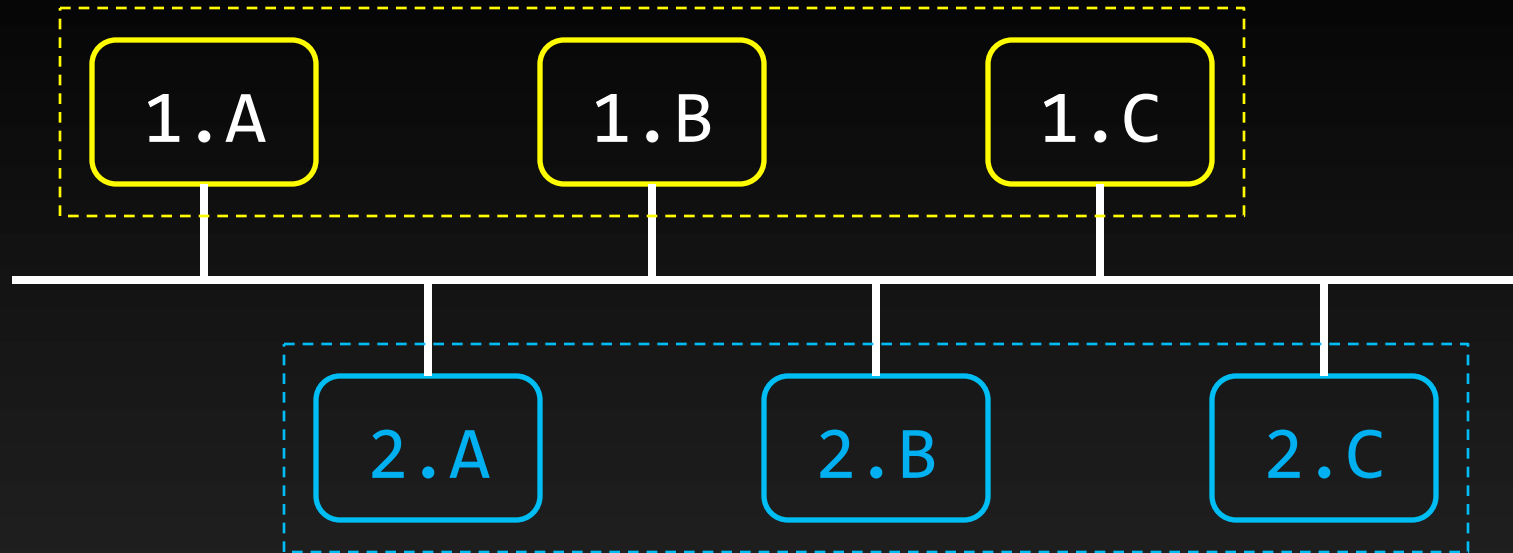
roll back

New?

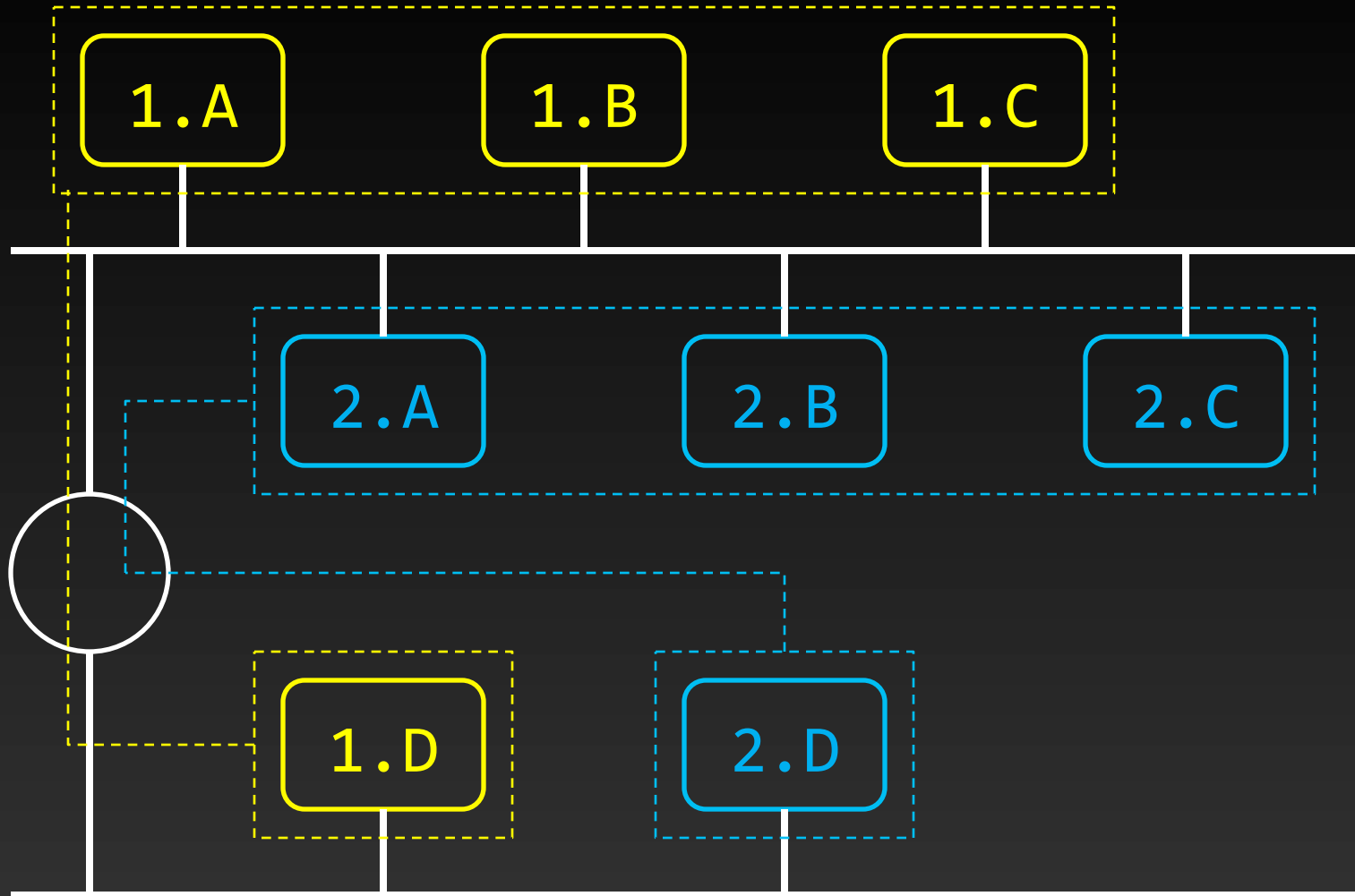
crude



less crude



less crude



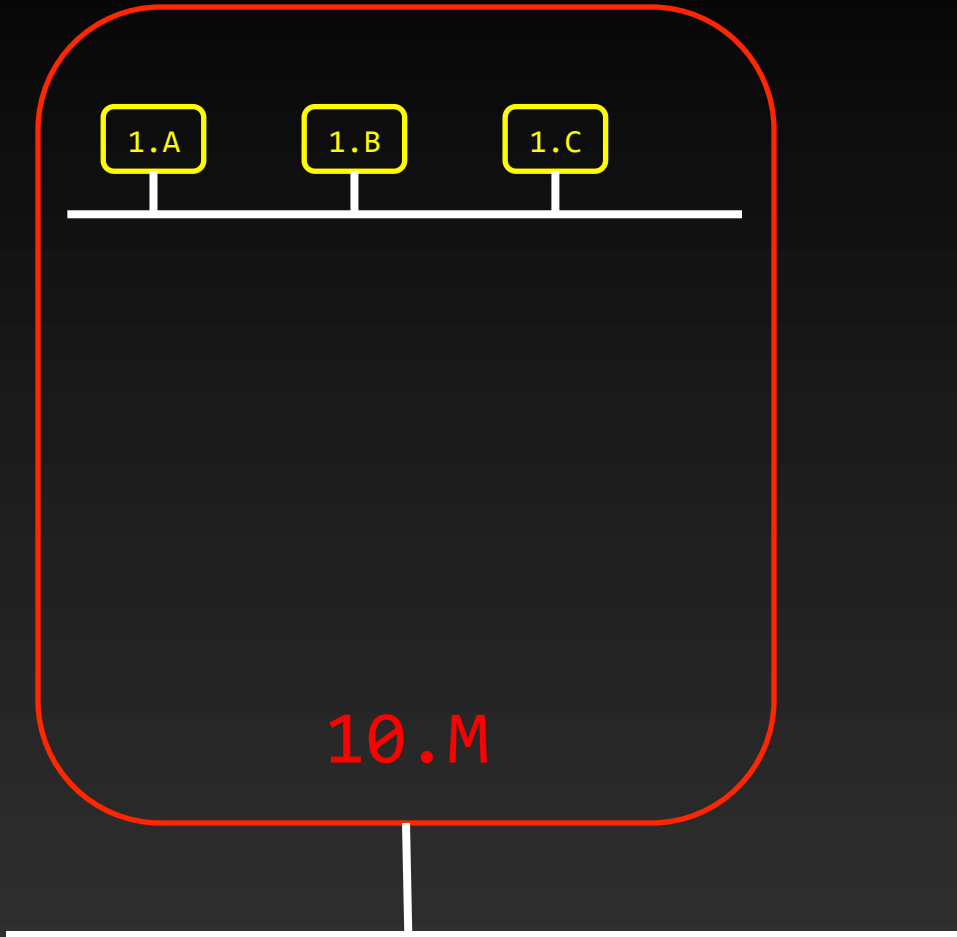
limitations

$$n < \infty$$

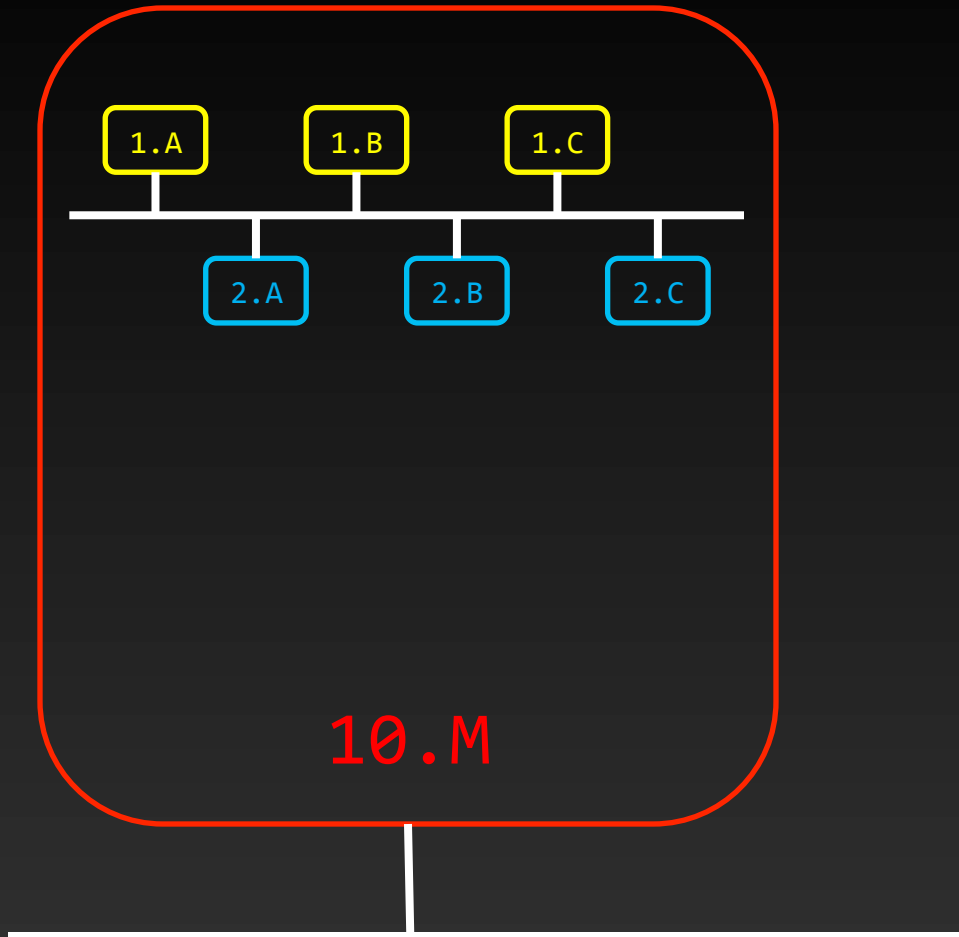
topology

static

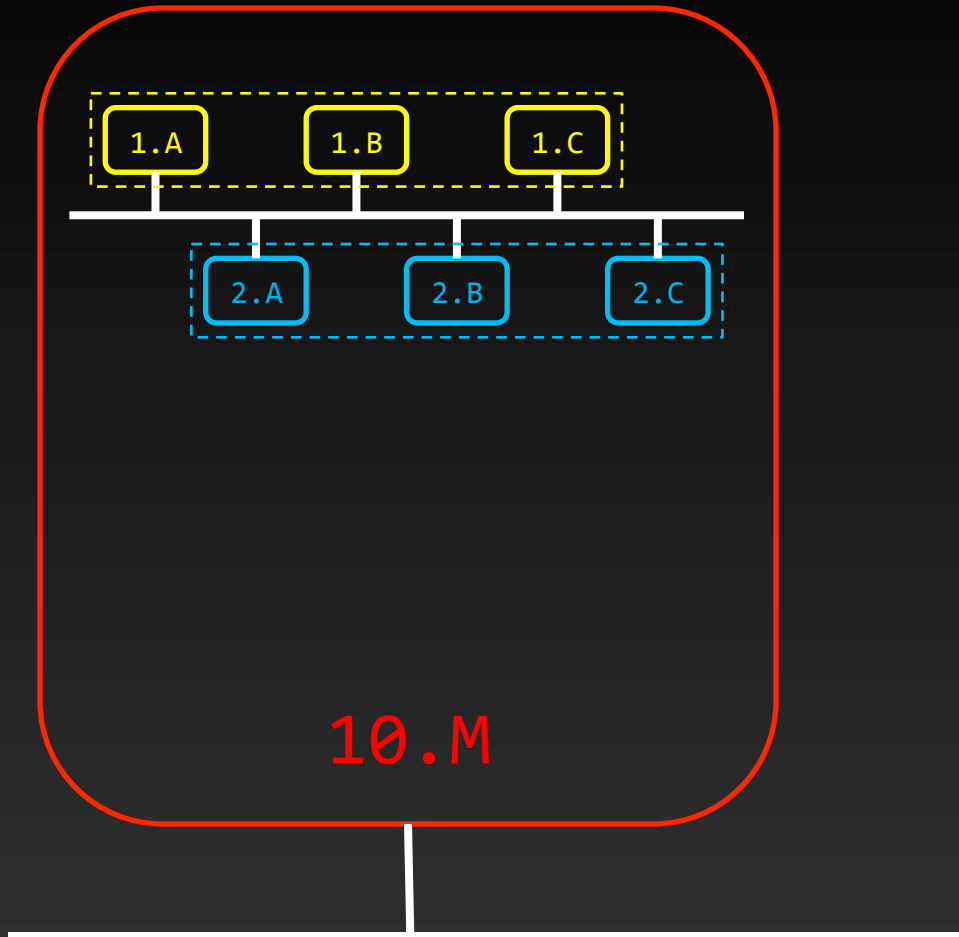
interesting



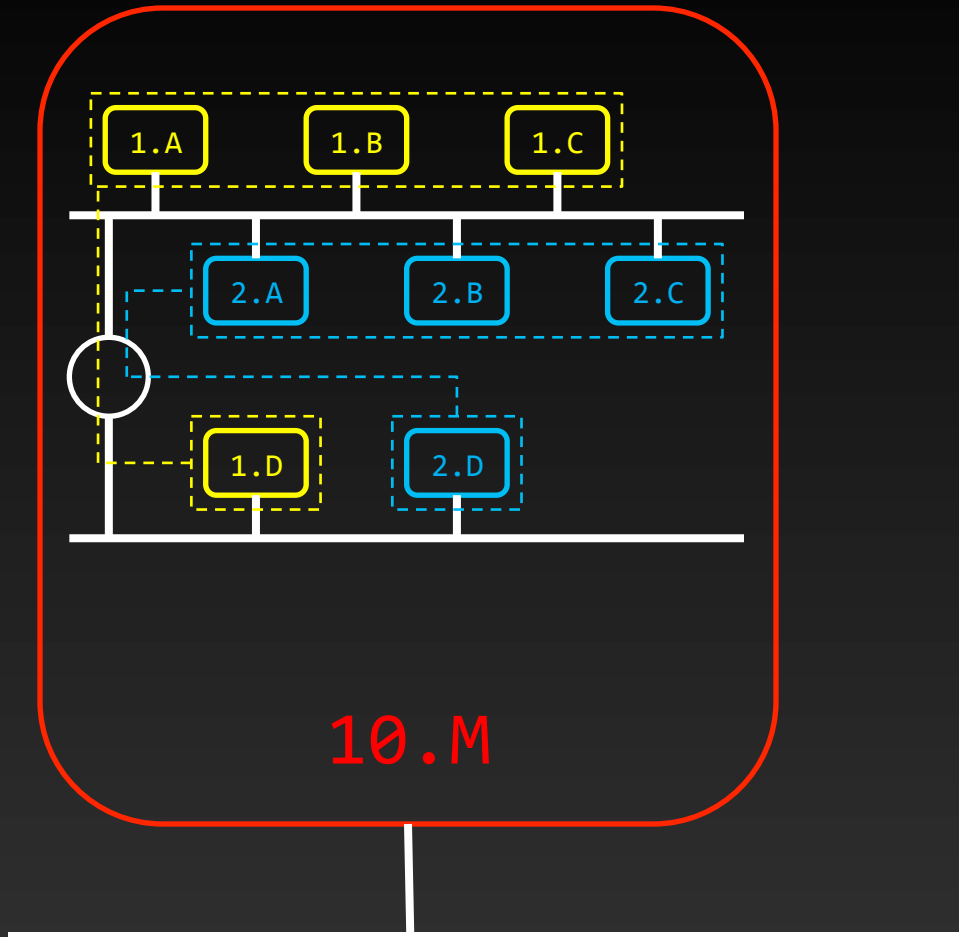
interesting +



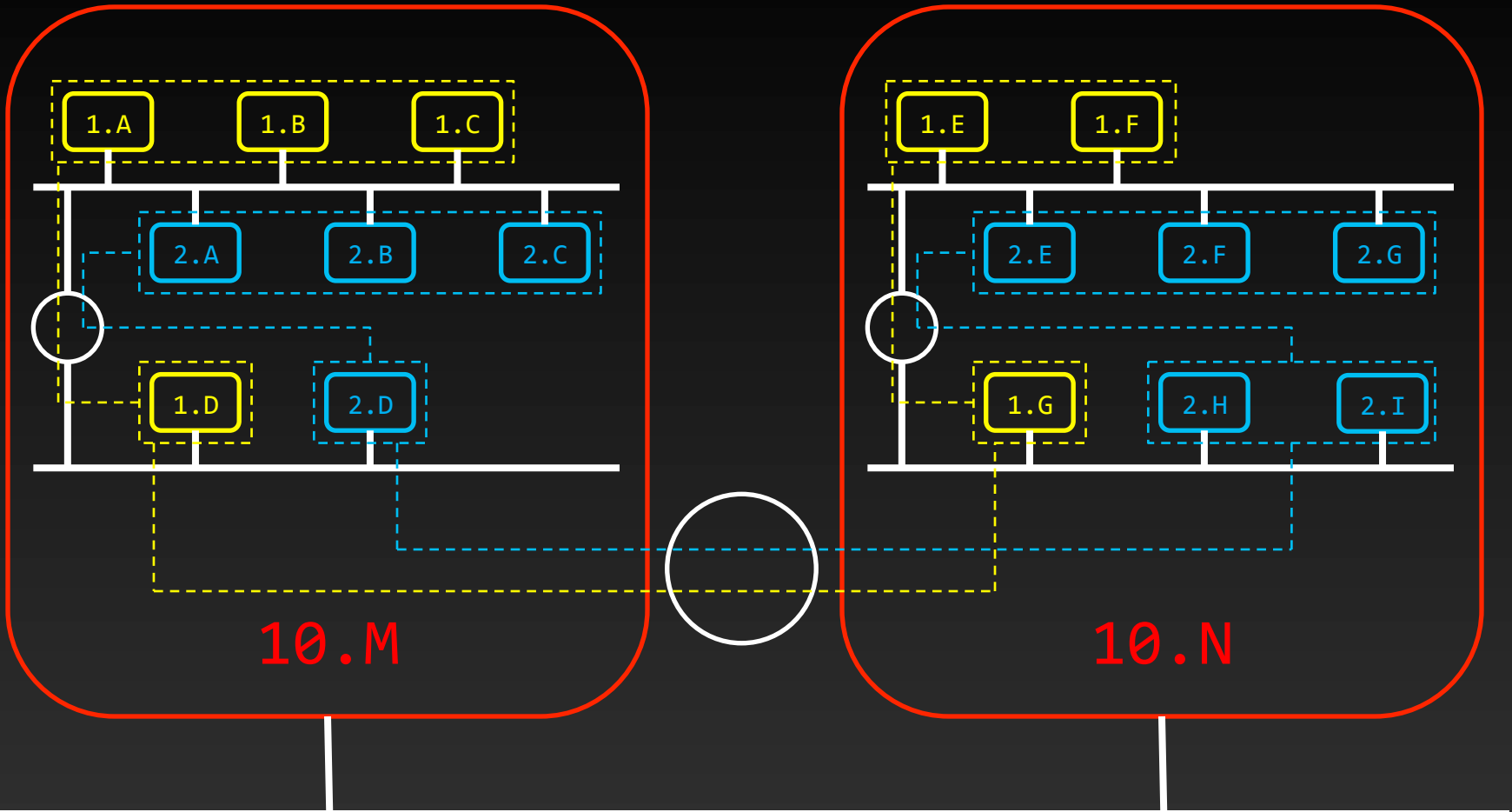
interesting +?



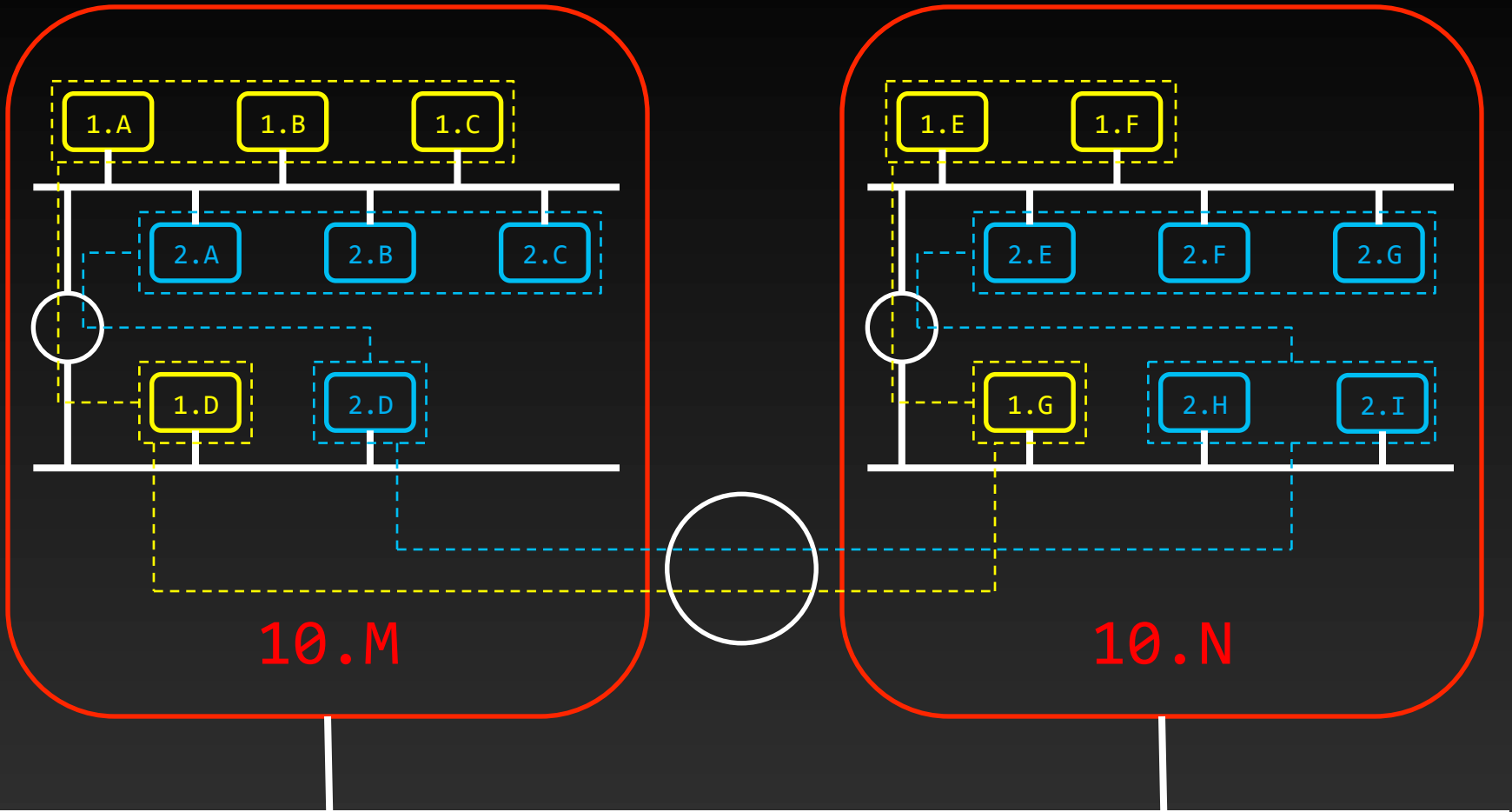
interesting +??



interesting +!!!



insane ;)



limitations

as before

+ not cloudable

operational abstractions
aren't useful

VM-1

VM-3

VM-2

VM-4

PM-1

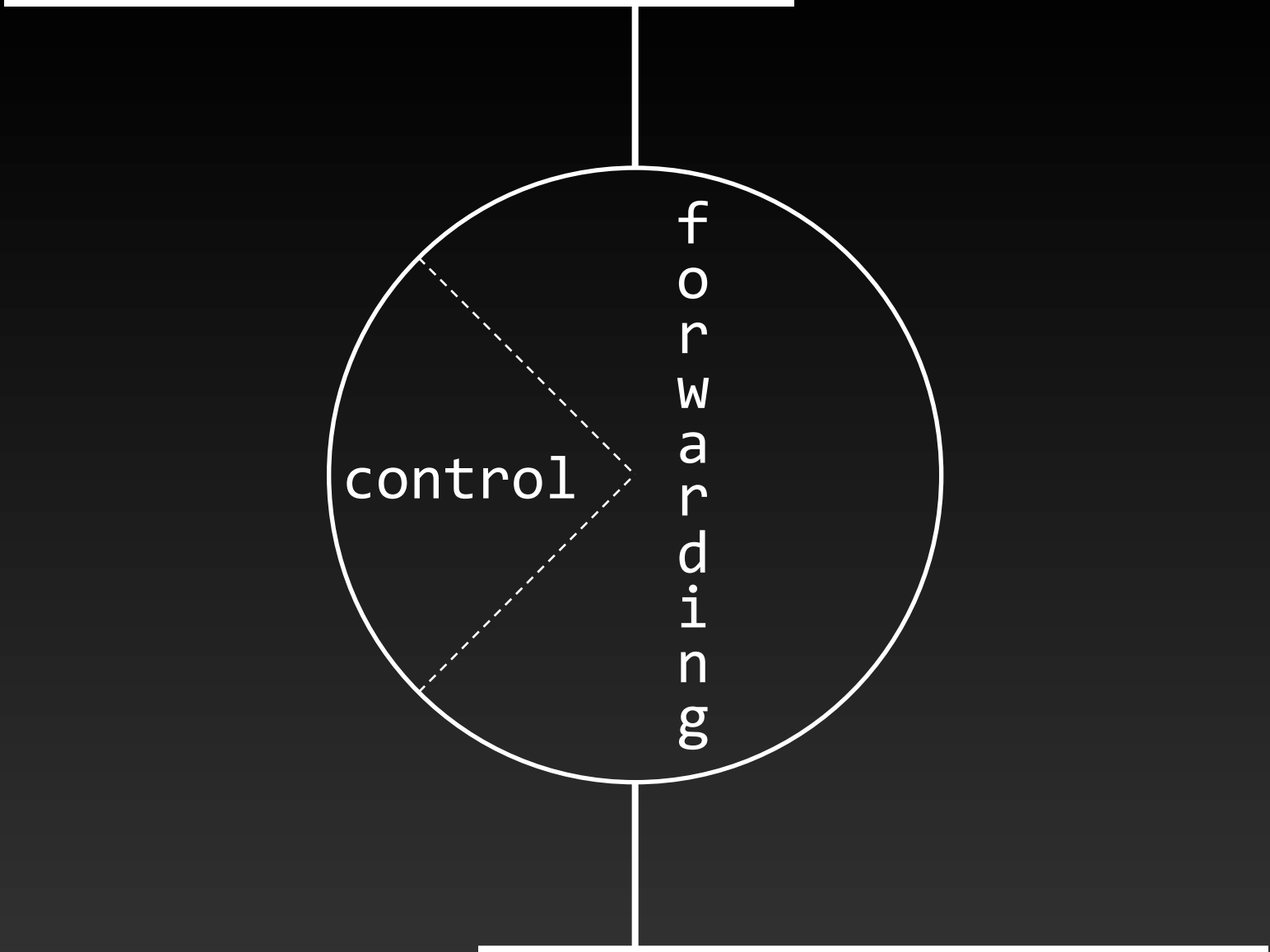
VM-1

VM-3

VM-2

MAC? IP? ACL? state?

PM-2



limitations

topology mandates/constraints

no overlapped addresses

sloooooow to change

+ not cloudable

requirements

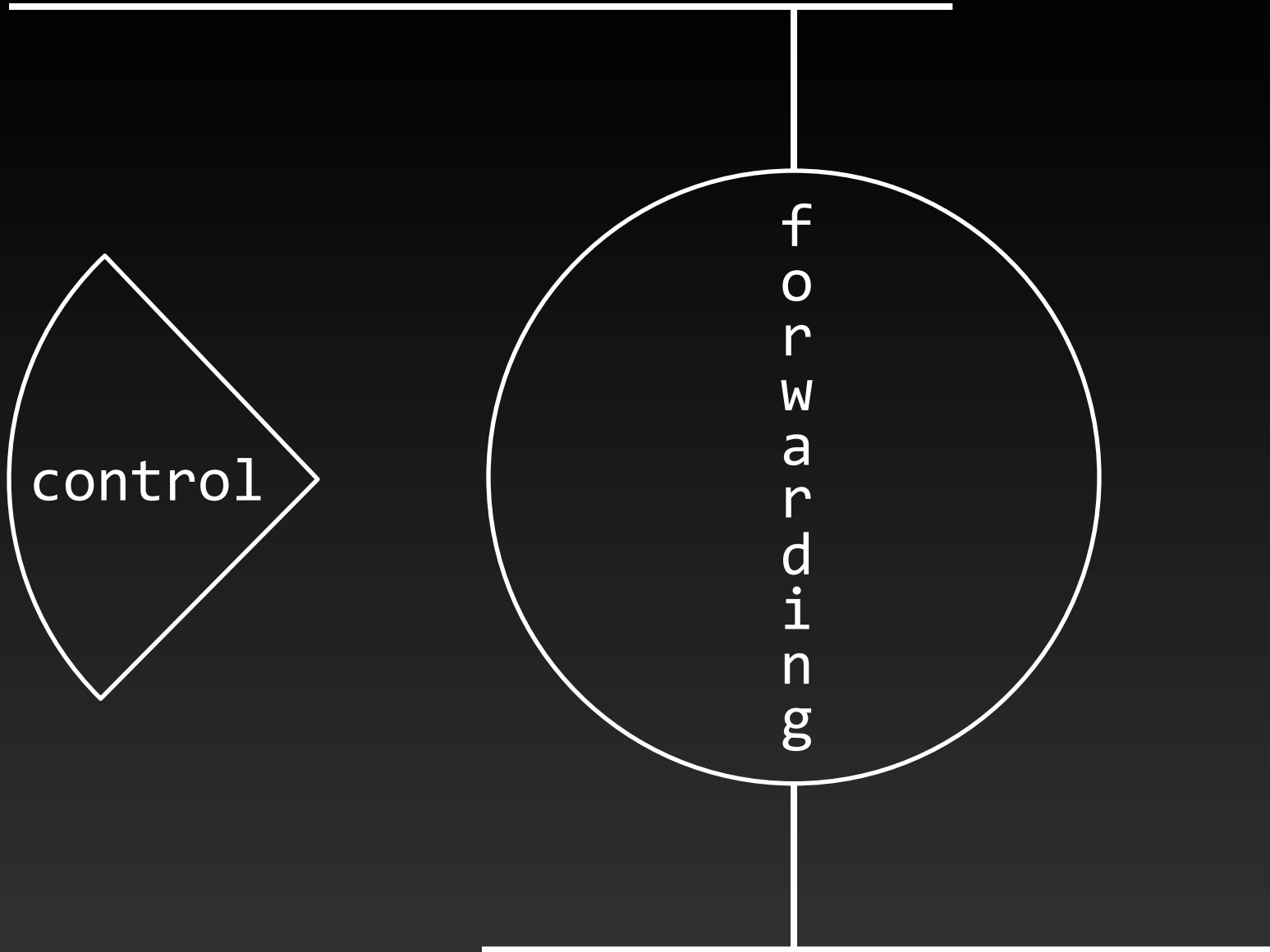
decouple V from P

V looks like P

V allows units of operation

Software Defined Networking (*)

* One popular, but not necessarily universal, definition

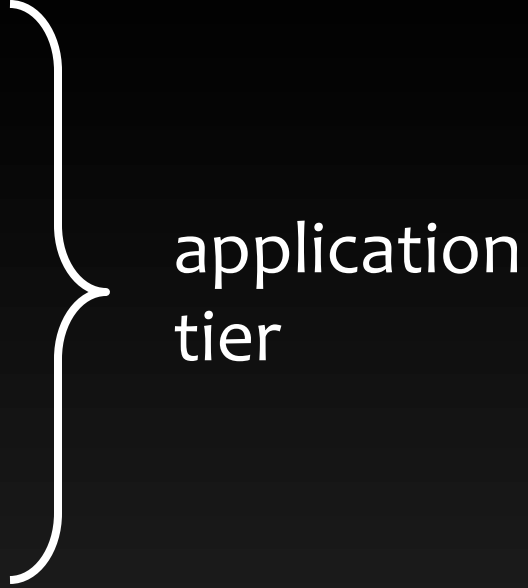
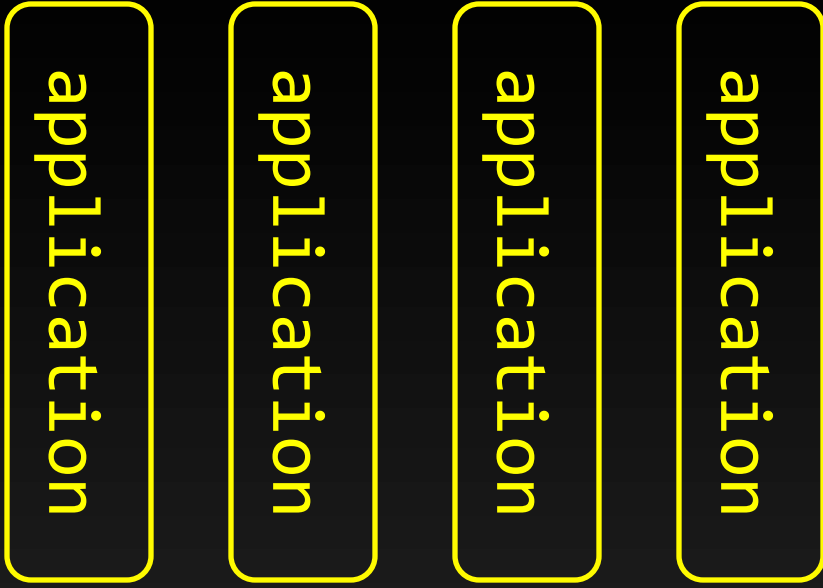


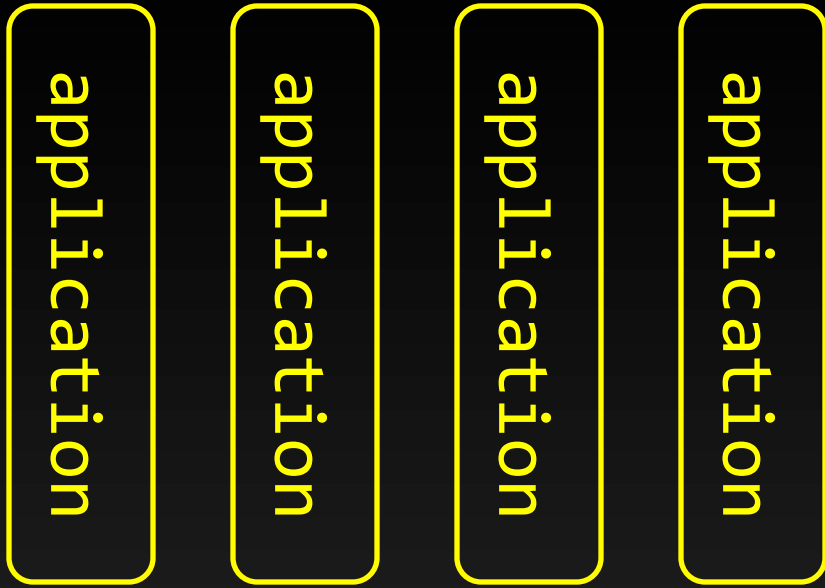
control

f
o
r
w
a
r
d
i
n
g

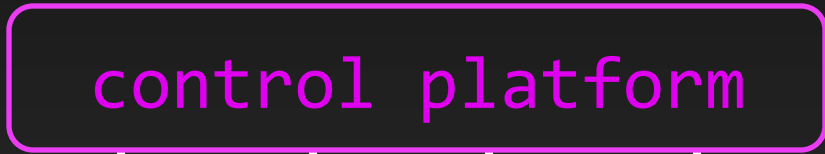
control

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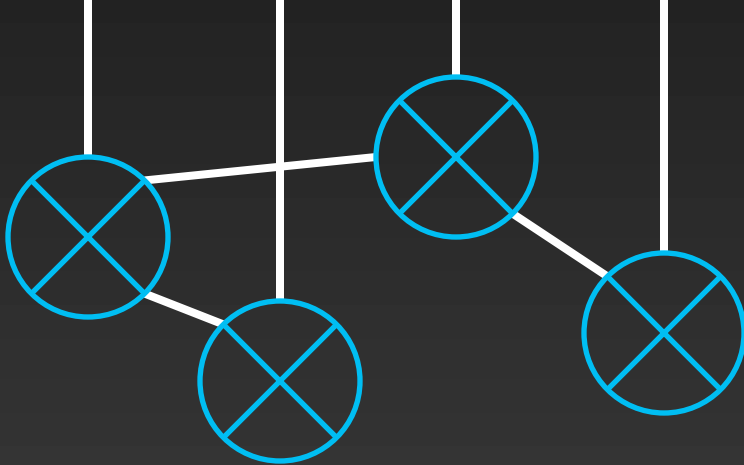




application tier



control plane



forwarding plane



OpenFlow

forwarding plane

relatively “dumb”

does what it's told

R.I.P.,

RIP

OSPF

IS-IS

&c.

control plane

centralized

end-to-end view
(not hop-by-hop)

programmable

naturally multitenant

maintains state

that's not?

virtual server

≠

virtual network

	Datapath	Consistency
Virtual server	CPU memory device I/O nanosecond operation	self-contained
Virtual network	address contexts	all-port knowledge N instances of N states consistency on all paths timely distribution

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The Virtual Network

decoupled from h/w

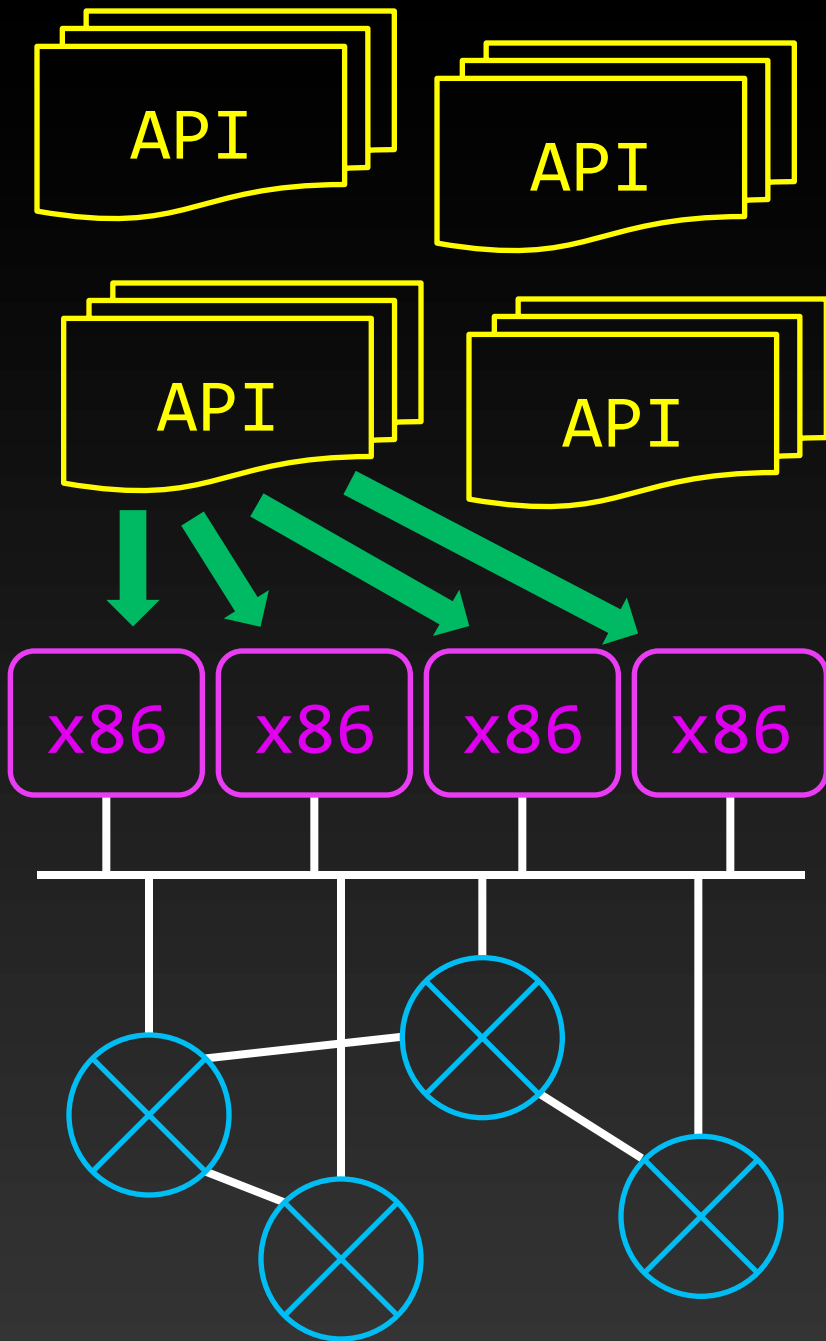
independent from others

delegated control

ephemeral

SDN (*) is a useful tool

* As defined previously



VXLAN
NVGRE
NVP
OTV
STT

control plane



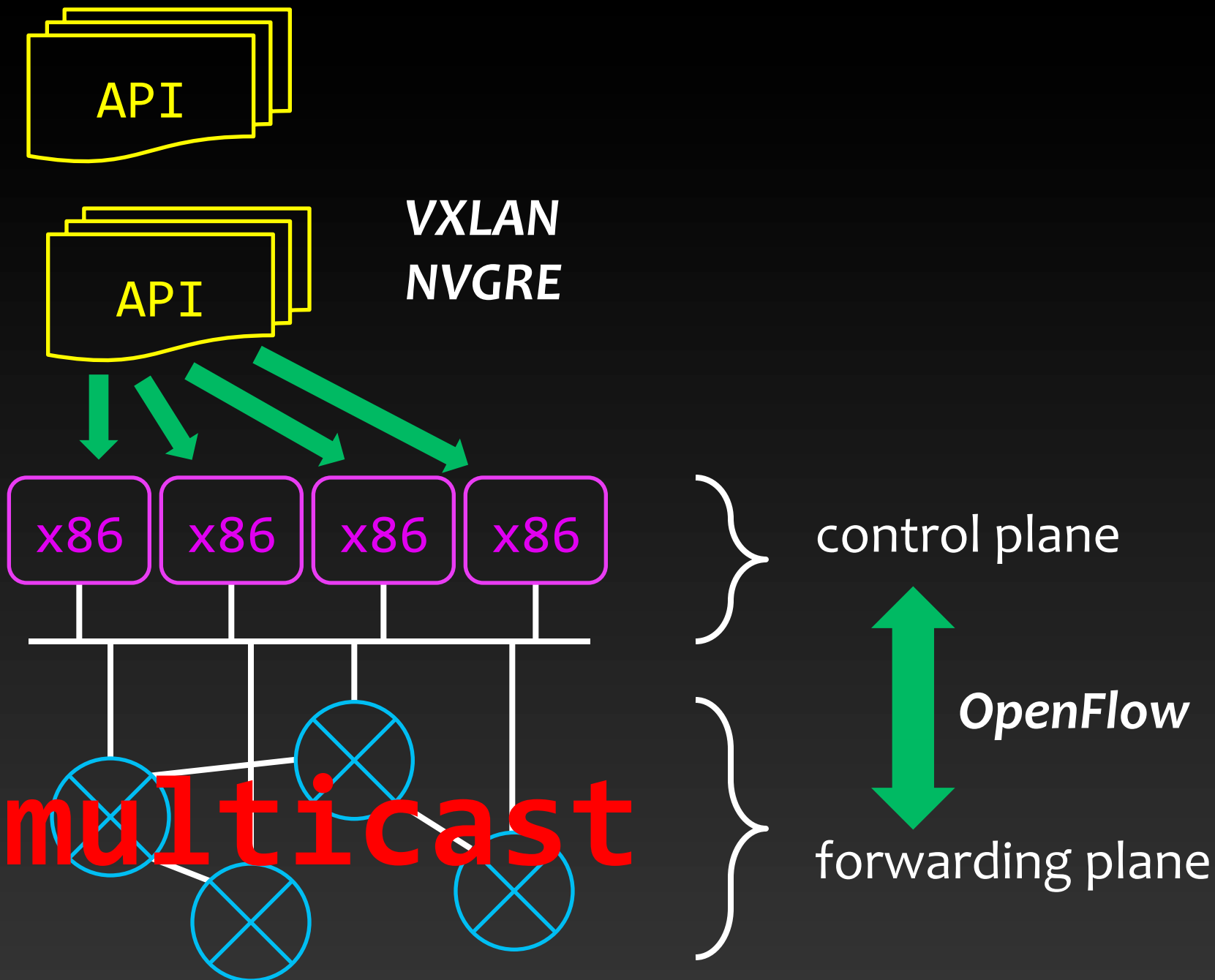
OpenFlow

forwarding plane



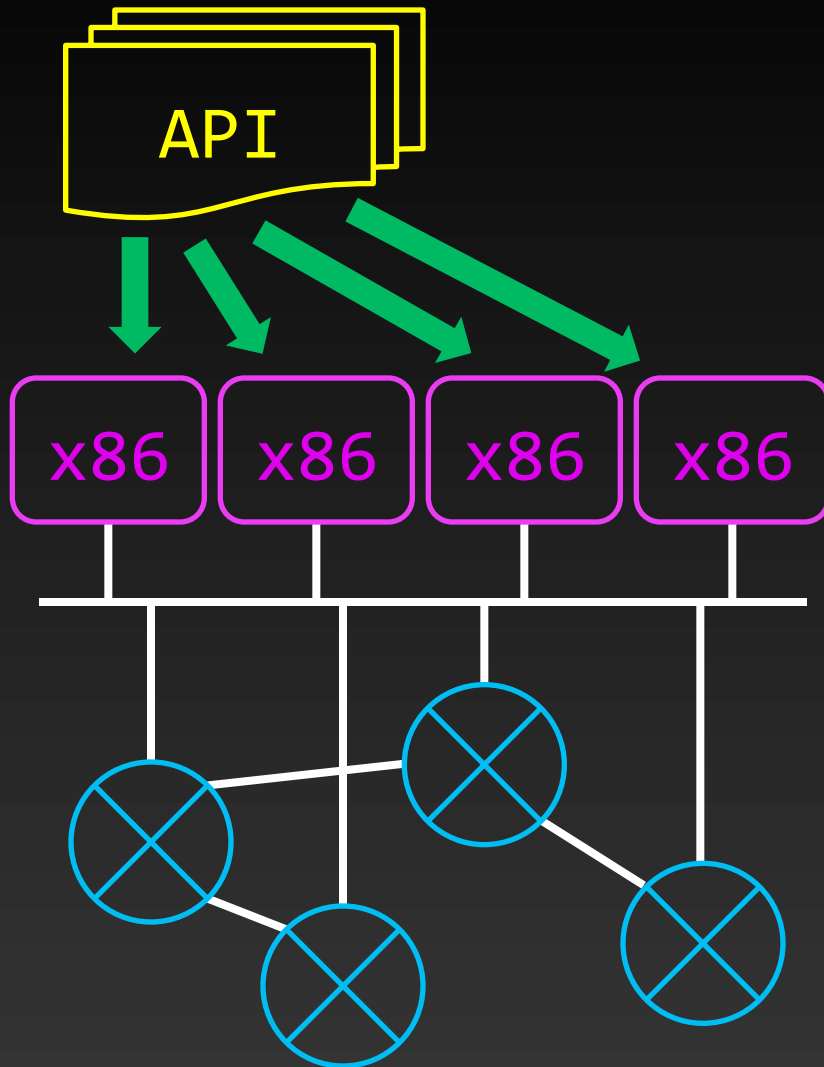
is SDN (*) a requirement?

* As defined previously





How does Alice talk to Bob?



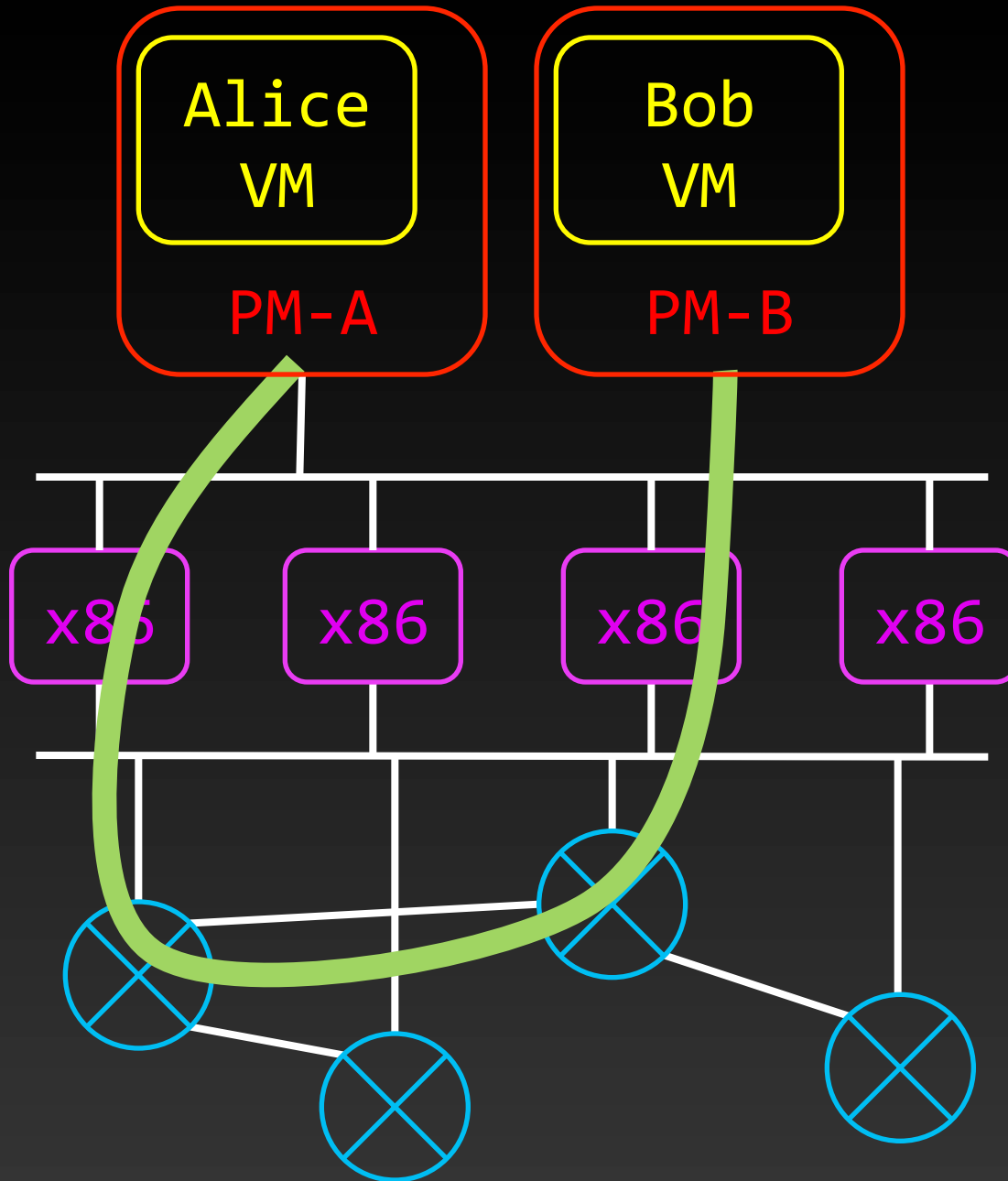
vAPI :

“I need a virtual L2-L3 network with these properties...”

vSWITCHes in x86 boxes determine optimal path

OPENFLOW:

“Hardware, plumb the following...”



E V I L



one flow, two VMs, separate hypervisors

	Throughput	Recv CPU	Send CPU
Linux bridge	9.3 Gbps	85%	75%
OVS bridge	9.4 Gbps	82%	70%
OVS-STT	9.5 Gbps	70%	70%
OVS-GRE	2.3 Gbps	75%	97%

aggregate, four VMs, two hypervisors

	Throughput	CPU
OVS bridge	18.4 Gbps	150%
OVS-STT	18.5 Gbps	120%
OVS-GRE	2.3 Gbps	150%

possibilities

security application

QoS application

WAN op application (*)

* *hard*: distributed cache and symbol vocabulary

on-demand VPN/C

infrastructure → code

network → code

*decoupled
and
delegated*

physical L2-L3



logical L2-L3
L4-L7 services

x86

x86, really?

complex → much CPU

FW/LB use CPU at flow start

optimized stacks ↑ performance

↑ upgrade certainty

distinct

security

forwarding

shaping

priority

...

outcomes

working multitenancy

isolated addressing

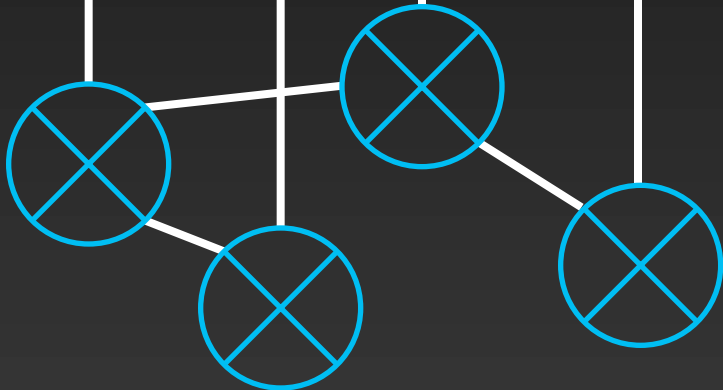
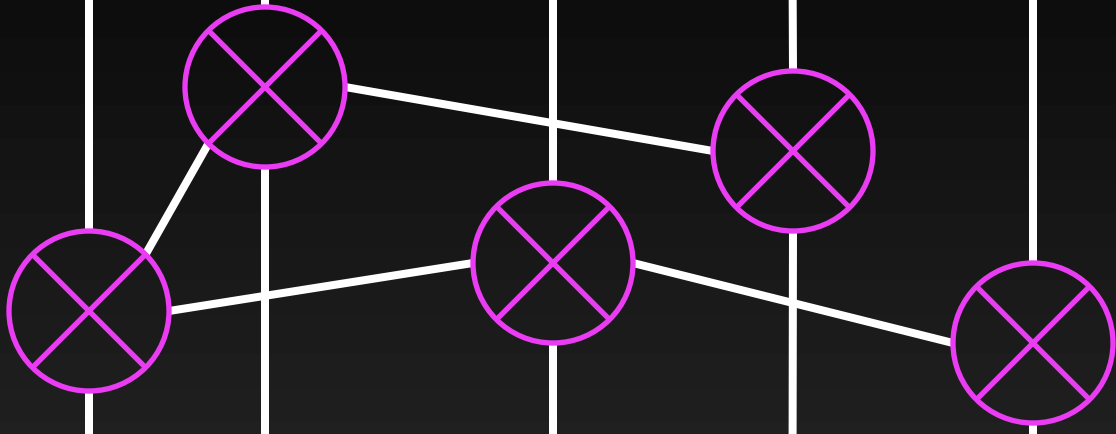
programmable

*independent
and
ephemeral*



“

”



virtual IP
virtual MAC

route my packets/frames
without collisions

move v-net
without changes

tear down when finished

separately alter
physical and virtual
topologies

consider:

on-demand HA/DR

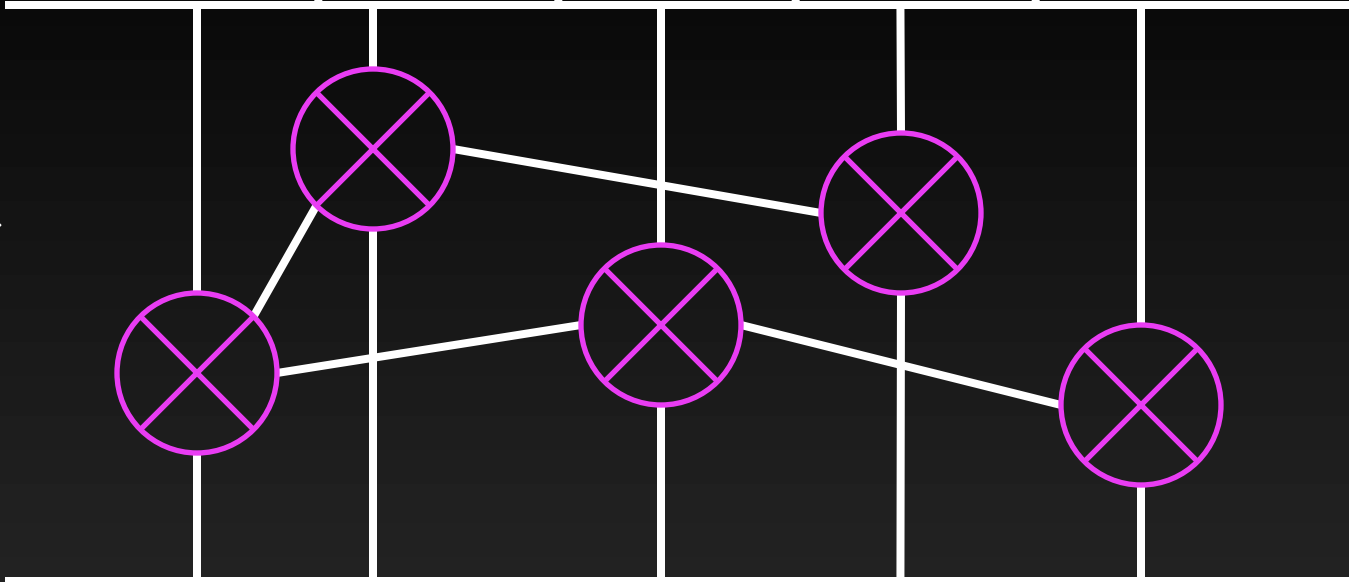
consider:
on-demand HDAR

SDN (*) manages state

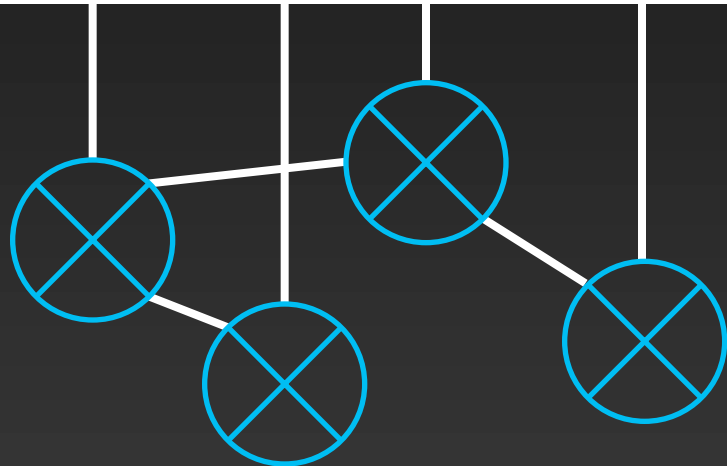
* As defined previously

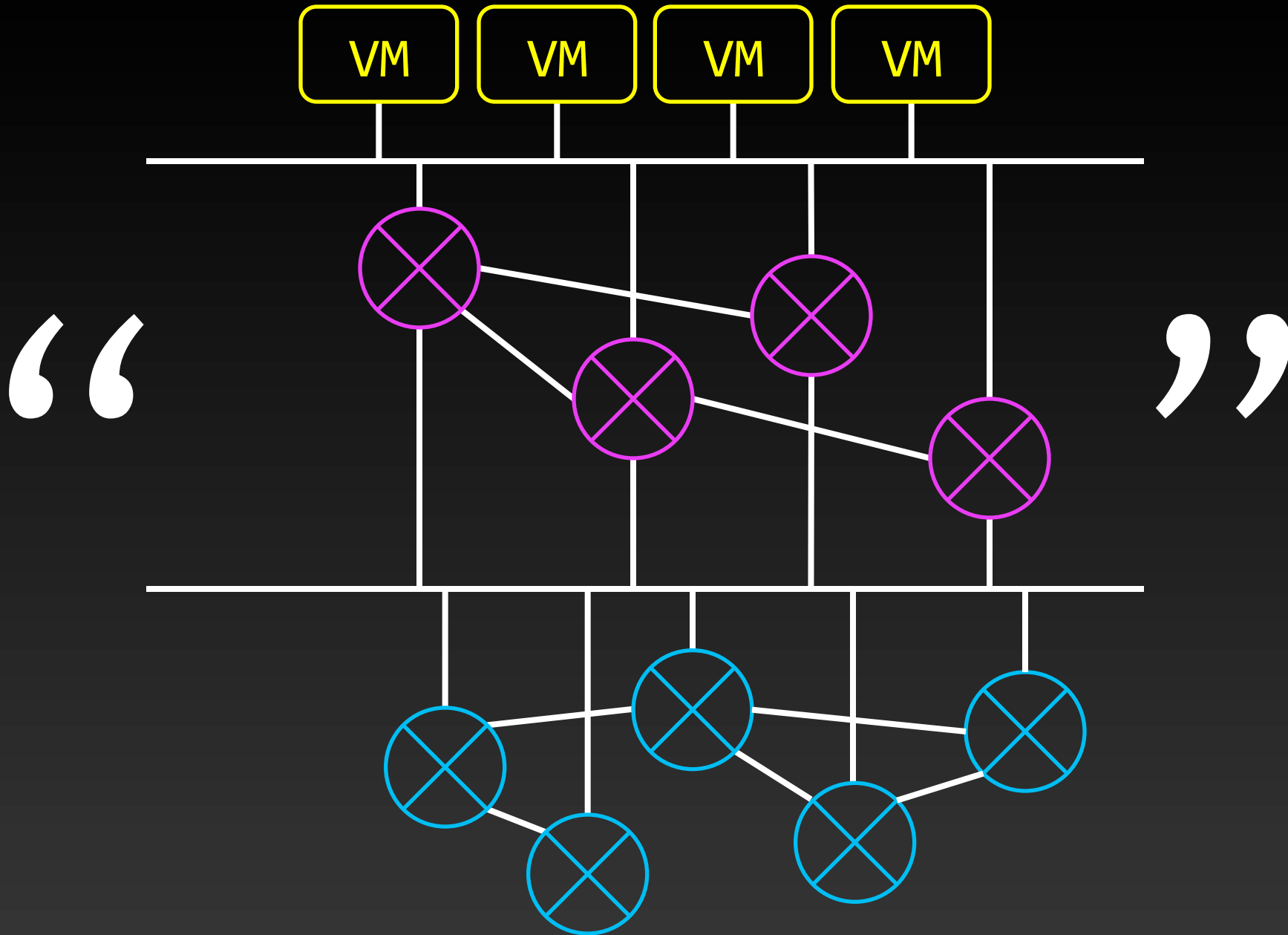


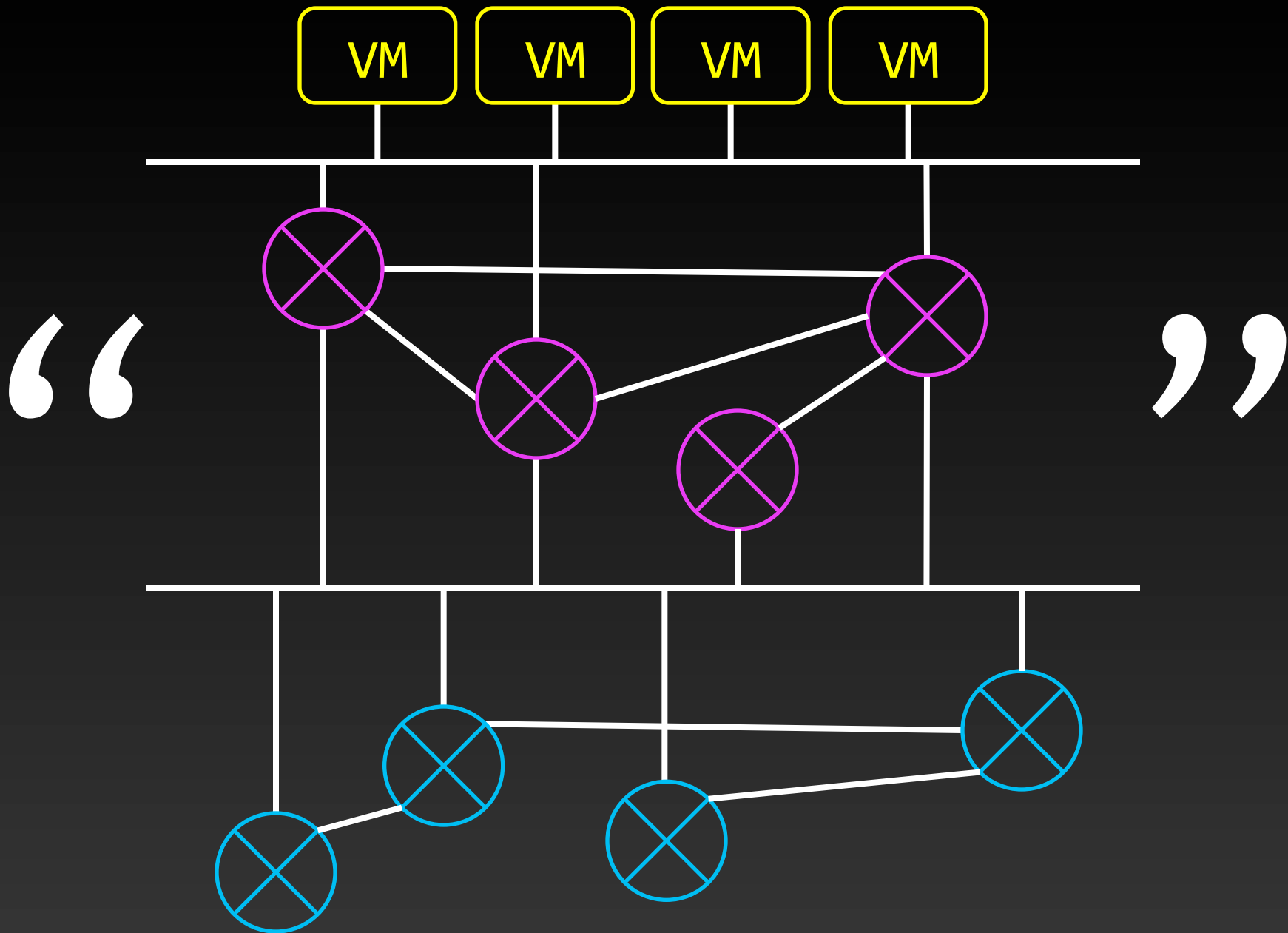
“



”







abstractional consistency

(mature orchestration?)

servers

=

disposable horsepower

networks

=

disposable pathways

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easy

familiar



point solution ideas

Tagging

Segmentation, not isolation
Same address in “both” worlds
Hardware has to understand
No mobility

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Address mapping	Like NAT: update address in place Multiplex large space into small: how? Virtual-to-virtual: physical “punch”

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Address mapping	Like NAT: update address in place Multiplex large space into small: how? Virtual-to-virtual: physical “punch”
Encapsulation	Or tunnels, or overlays (sigh) Worlds can be totally distinct Different forwarding for V and P Strong isolation: no V on P w/o bridge



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programmability
and cloudability



hard

scary



innovative advancements

Resources

networkheresy.com

packetpushers.net

blog.ioshints.com

sdncentral.com

Thanks for coming!

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